Comparison Analysis of Graph Theory Algorithms for Shortest Path Problem

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Abstract- The Sumba region, Indonesia, is known for its extraordinary natural beauty and unique cultural richness. There are 19 interesting tourist attractions spread throughout the area, but tourists often face difficulties in planning efficient visiting routes. From this case, it can be solved by applying graph theory in terms of searching for the shortest distance which is completed using the shortest path search algorithm. Then these 19 tourist objects are used to build a weighted graph, where the nodes represent the tourist objects and the edges of the graph describe the distance or travel time between these objects. Therefore, this research aims to compare the shortest path search algorithm with parameters to compare the shortest distance results, algorithm complexity and execution time for tourism in the Sumba area. The results of this research involve a comparison of several shortest path search algorithms, with the aim of finding the shortest distance results, algorithm complexity, and execution time for tourism in the Sumba area. Based on the test results of the five algorithms with the parameters that have been prepared, and the findings show that each algorithm has its own characteristics, the results are as follows: Dijkstra's algorithm can be used to calculate the shortest route for single-source and single-destination types. This resembles the Bellman-Ford algorithm, only the Bellman-Ford algorithm can be used simultaneously on graphs that have negative weight values. Meanwhile, the Floyd-Warshall algorithm is suitable for use on the all-pairs type. Then, the Johnson Algorithm can be used to determine the shortest path from all pairs of paths where the destination node is located in the graph. Finally, the Ant Colony algorithm to compute from a node to each pair of destination nodes.

Keywords— Dijkstra Algorithm, Bellman-Ford Algorithm, Floyd-Warshall Algorithm, Johnson Algorithm, Ant Colony Algorithm

I. INTRODUCTION

Graph theory is one of the topics in the computer field that is studied and is a branch of discrete mathematics [1] that studies related graphs, namely the relationship between one object and another object, where the connection consists of a set of vertices/points (vertex/nodes) and sides/lines (edges) that connect one point to another point. Graph theory is used in studies related to relations or relationships between objects that are implemented to solve various problem models, one of which is the optimal path or route search problem [2-3] so that it can minimize costs [4] and time efficiency [5], optimization scheduling [6-8], communication network modeling [9], and other issues [9]. Searching for the shortest path is the process of finding a path between two vertices, namely from the source node to the destination node in a graph by minimizing the number of weights so that the path with the smallest weight is obtained.

The Sumba region is a tourist destination rich in attractive natural and cultural tourist attractions, and 19 attractive tourist attractions are distributed throughout the region. However, for many tourists, the challenge often faced is to plan efficient visits to these various tourist attractions on a single trip. This is becoming increasingly important as their time and resources are limited. To address this problem, this study applied the implementation of graph theory in terms of the shortest distance search between tourist attractions in Sumba area through testing of several algorithms that have often been used to solve related problems. With this effort, it is also possible for each algorithm to perform efficiently on what kind of graph and path models, as well as what the advantages and disadvantages of each algorithm are. The hope can also be a recommendation for readers to know the most efficient algorithms for researching or solving the same problem.

Implementation of graph theory in the case of finding the shortest distance can be solved using the shortest path search algorithm, which is classified into a single source and multisource [10] where the single source algorithm is the shortest path search algorithm from one source node to various other destination nodes. These algorithms are among the Greedy [11], Dijkstra, and Bellman-Ford algorithms. Meanwhile, multi-source algorithms look for the shortest path by calculating all pairs of vertices in a graph, including the Floyd Warshall, Ant Colony, and Johnson algorithms. Several algorithms have made it possible for tourists to visit as many tourist attractions as possible in one visit, while minimizing travel time.

Some research using the Greedy algorithm includes determining the best tour packages for tour communicatists [12] and determining the fastest routes for public transportation [13]. The application of Dijkstra's algorithm

p-ISSN 2301-7988, e-ISSN 2581-0588 DOI : 10.32736/sisfokom.v12i3.1756, Copyright ©2023 Submitted : May 24,2023, Revised : August 8, 2023, Accepted : Oktober 1, 2023, Published : November 6, 2023 includes determining the shortest path [14-21], that there is the application of the Bellman-Ford algorithm in finding the best path including the network [22], a comparison of the Bellman-Ford algorithm with other algorithms such as Prim's algorithm [23], Dijkstra [24-27]. Regarding the application of the Floyd-Warshall algorithm in finding the shortest route, they include [16][28]. The application of other shortest route search algorithms such as Ant Colony includes [1][20-21][29-30]. Meanwhile, the application of Johnson's algorithm in finding the shortest path [6]. There are also other studies that make comparisons between the Dijkstra, Bellman-Ford, and Floyd-Warshall algorithms [10].

The shortest path search algorithms that have been mentioned have their own advantages and disadvantages which can be measured through several parameters including, the results of the shortest path or route given [24-25], negative sides [26], negative cycles [26][10], the memory allocation used [10], the complexity of the algorithm [24], as well as the execution time required for an algorithm [26][10].

This study has the primary purpose of testing and comparing from five different shortest path search algorithms in optimizing the travel of tourists in the Sumba region. This area is a focal point because it presents its own challenges in determining tourist attractions to visit, along with the condition of road infrastructure that still suffers a lot of damage. Therefore, this study will focus on a number of relevant parameters, including shortest distance results, algorithm complexity, and execution time.

II. METHODOLOGY

The data used in this research is about the distance data between tourist objects in Southwest Sumba and West Sumba. The data is taken from google maps and consists the distance between TMC Airport to Kita Beach, and also to Kabisu Lobo Oro Site, Lendongara Hill, Waikelo Beach, Kawona Beach, Karakar Indah Beach, Waikuri Lagoon, Mandorak Beach, Tanjung Karoso Beach, Tanjung Karoso Resort, Rotenggaro Beach, Bawana Beach, Watu Malandong Beach, Sumba Culture House, Praijing Village, Lailiang Beach, Rua Beach, and Mata Yangu Waterfall. The route data was taken with the consideration of good road access, paved and two-way without any damage.

The following is an explanation of the initialization of the dataset to be used. To facilitate the writing of data tables, each location name will be represented by numbering as shown in Table 1 of initializing variables. Then, continue with Fig. 1 that displays information about the distance values in each row and column that have been adjusted by initializing the location variables and distance values using units of kilometers.

Variable	Sumba	Variable	Sumba			
Initialization	Regional	Initialization	Regional			
	Destination		Destination			
1	TMC Airport	11	Tanjung Karoso			
			Resort			
2	Kita Beach	12	Rotenggaro			
			Beach			
3	Kabisu Lobo	13	Bawana Beach			
	Oro Site					
4	Lendongara Hill	14	Watu			
			Malandong			
			Beach			
5	Waikelo Beach	15	Sumba Culture			
			House			
6	Kawona Beach	16	Praijing Village			
7	Karakat Indah	17	Lailiang Beach			
,	Beach	17	Luniung Deuen			
8	Waikuri Lagoon	18	Rua Beach			
9	Mandorak	19	Mata Yangu			
	Beach		Waterfall			
10	Tanjung Karoso					
	Beach					

FROM (TO	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19
1	0	19	9	10	8,3	8,7	22	42	40	51	43	47	58	63	6,5	43	66	59	65
2	19	0	20	13	22	26	40	59	57	67	59	63	75	79	23	43	68	61	67
3	9	20	0	11	12	16	30	49	47	57	49	53	64	69	13	49	72	65	71
4	10	13	11	0	13	17	31	50	49	58	50	54	66	70	14	50	73	66	73
5	8.3	22	12	13	0	5.7	19	40	38	54	46	50	62	66	11	47	70	63	70
6	8.7	26	16	17	5.7	0	14	34	32	49	42	45	56	61	12	49	72	64	71
7	22	40	30	31	19	14	0	20	18	35	28	31	43	47	17	57	80	73	79
8	42	59	49	50	40	34	20	0	2	11	9.1	20	31	36	38	78	101	94	100
9	40	57	47	49	38	32	18	2	0	9	7.1	18	29	34	36	76	99	92	98
10	51	67	57	58	54	49	35	11	9	0	5.7	14	26	30	43	83	106	81	105
11	43	59	49	50	46	42	28	9.1	7.1	5.7	0	18	30	35	37	77	101	86	100
12	47	63	53	54	50	45	31	20	18	14	18	0	19	24	41	81	95	75	104
13	58	75	64	66	62	56	43	31	29	26	30	19	0	7.9	53	72	79	59	95
14	63	79	69	70	66	61	47	36	34	30	35	24	7.9	0	57	67	74	54	90
15	6.5	23	13	14	11	12	17	38	36	43	37	41	53	57	0	41	64	57	64
16	43	43	49	50	47	49	57	78	76	83	77	81	72	67	41	0	19	21	24
17	66	68	72	73	70	72	80	101	99	106	101	95	79	74	64	19	0	20	51
18	59	61	65	66	63	64	73	94	92	81	86	75	59	54	57	21	20	0	43
19	65	67	71	73	70	71	79	100	98	105	100	104	95	90	64	24	51	43	0

Fig. 1. Distance Data between Tourist Objects

The data above are initial data showing information on the distance values between one point to each other obtained at the data collection stage. The dataset will be used in the testing process, as well as to see if a particular algorithm can find a solution so that a point can have a smaller distance value while also finding the shortest route between the starting point and the destination point. The data consists of 19 location points with each distance value written with a unit of kilometers. This study also aims to optimize the travel of tourists in the Sumba region. This study wanted to contribute to facilitating the determination of tourist attractions to be visited, especially when talking about road infrastructure conditions and sometimes there is a potential route mismatch that certain applications can produce when they want to know the route on their way to a location.

The algorithms used in this study include Dijkstra, Bellman-Ford, Floyd-Warshall, Johnson, and Ant Colony algorithms. Dijkstra's algorithm is an algorithm that is used to find the shortest distance or route from the initial node to the final node in a weighted graph. Dijkstra's algorithm is that it is the most important and useful algorithm for optimal solutions in a large class of shortest path problems [14]. Bellman-Ford algorithm can also be used to solve problems related to the shortest path problem where this algorithm can work even though there are negative edge weight values. This algorithm itself is a refinement of Dijkstra's algorithm. Bellman-Ford algorithm offers a dynamic solution for all nodes in a graph to determine the minimum route with edge weights that can be negative, but still does not contain negative cycles [23]. Floyd-Warshall algorithm is a dynamic algorithm and is often used to determine the shortest route between all points on a directed graph without negative cycles. This algorithm is presented to find solutions to the shortest-path problem between certain fixed nodes and other related nodes [16]. Johnson's algorithm is also one of the commonly used algorithms to solve the shortest route problem. This algorithm is a combination of the Bellman-Ford and Dijkstra algorithms. Johnson's algorithm, among other things, can be used on negative weighted graphs. [6] In addition, Johnson's algorithm is also an appropriate solution method for solving scheduling problems such as the research by M. Okwu and I. Emovon [7] and by M. Redi and M. Ikram [8]. Ant Colony Optimization (ACO) algorithm is an optimization algorithm that uses probabilistic techniques and is used to solve computational problems and find optimal paths. The Ant Colony Optimization strategy will choose the shortest path based on the path most frequently traveled by ants, using mechanisms that mimic behavior or social strategies that exist in nature [31]. In testing the five algorithms, it is also assessed based on the type of shortestpath, regarding which algorithm is suitable for use in finding solutions based on a particular type of shortest-path. There are several types such as single-pair, then single-source, single-destination, and all-pairs shortest path problem.

The parameter used as an indicator of the assessment of an algorithm being tested. The first is related to the calculation of the shortest distance, then also related to complexity, execution time, and the final result of the distance that being traveled. Calculation of the shortest distance is to find the path between the initial node and the final node so that the number of edges with minimum weight is found. For example, that a node can be connected directly to another node through one edge or a series of edges. So this research has observed visually, that there may be shorter 'distances' between some vertices and others in the graph [17]. Complexity is a parameter that provides information regarding how complicated aspects of the algorithm used are. Complexity here actually has a broader definition, in the sense that this parameter can cover several aspects such as the complexity of memory usage, space, running time [18][19][20], and so on. Execution time is used to find out how efficient and effective an algorithm is in terms of time to process data. In research by Abusalim, et al. [27]Studying that the number of nodes that make up the graph has an influence on the fast or slow running time when executing certain algorithms. Then, for the final result is to talk about how the output is generated in the program.

The research was carried out through several stages, such as receiving input in the form of data used, then proceeding with the determination of the starting point and destination point used during the process. Implementation was carried out on five algorithms and through these implementations, followed by an analysis based on the parameters that have been compiled to obtain results. The following Fig. 2 is to describe the flow of the research through a flowchart.



Fig. 2. Distance Data between Tourist Objects

Based on the flowchart, the first step is to collect the data. In this case, the dataset that will be used in the research. The stage continues with a literature review regarding each algorithm that will be used as material in testing. And also from this dataset, the starting point and destination point are determined as a starting point for starting the test. After determining the point of reference, the implementation is

p-ISSN 2301-7988, e-ISSN 2581-0588 DOI : 10.32736/sisfokom.v12i3.1756, Copyright ©2023 Submitted : May 24,2023, Revised : August 8, 2023, Accepted : Oktober 1, 2023, Published : November 6, 2023 carried out with five algorithms with the help of the Python program. From the tests carried out, then the results and comparisons are analyzed referring to the parameters that have been compiled, related to the calculation of the shortest distance, complexity, execution time, and the final result of the distance traveled. Then the stage is continued by determining the results and conclusions obtained through the tests that have been carried out.

III. RESULT AND ANALYSIS

The results and analysis of the shortest-path search calculation use 19 data from the Sumba Region nodes with five algorithms, including the Dijkstra, Bellman-Ford, Floyd-Warshall, Johnson, and Ant Colony algorithms. From these algorithms, modeling is implemented using Python programming language with Spyder application (Python 3.7). For the explanation as follows.

A. Dijkstra's Algorithm

1) Calculation of the shortest distance

The steps for calculating Dijkstra's algorithm use an approach in determining the shortest path starting from a TMC Airport node to several other destination nodes in a graph, including the following :

(a) The first step is to mark all the nodes to be visited,

(b) Mark the selected initial node with current distance 0, and other nodes with infinity " ∞ ",

(c) Then update the initial node as the current node,

(d) For the current node, analyze all unvisited neighbor nodes and measure their distance by adding the current distance from the current node to the edge weight connecting the neighbor node and the current node,

(e) Compare the distance measured recently with the currently assigned distance to the neighboring nodes and take that as the new current distance from the neighboring node,

(f) After that, consider all unvisited neighbors of the current node, mark the current node as visited,

(g) If the marked destination node is visited then stops, then the algorithm has ended. If not, select the unvisited node marked with the closest distance, fix it as the current new node, and repeat the process again from step (d).

2) Complexity

The complexity when running Dijkstra's algorithm program in Python is 70 lines.

3) Execution Time

Dijkstra's Algorithm is carried out 5 times to run the Python language program, so that it can better determine the amount of time the program is running, and can take the average time. The following Table II is for the test results related to the execution time of the program.

Execution	Time
First Execution	0.001322299999982221 seconds
Second Execution	0.000320000000020964 seconds

Third Execution	0.0003347000000335176 seconds
Fourth Execution	0.0004951999999960321 seconds
Fifth Execution	0.000729599999996639 seconds
Average Time	0.000640359999999873252 seconds

execution. For the first the program takes 0.0013222999999982221 seconds to run the algorithm to the stage of displaying the results obtained, which in this case is route information. This is also the same for the second execution up to the fifth execution stage. From the Table II, it can be seen that the time used to run the algorithm can be said to be very fast and does not have a significant difference in each process. The average execution time obtained was 0.000640359999999873252 seconds, meaning that programs that implement the Dijkstra's Algorithm can be executed by the console lightly and quickly.

4) Final Result

The results of the shortest distance that being traveled are calculated from TMC Airport to all destinations in the Sumba Region, with the results of the Table III data as follows.

No	Distance	Path Travarsad	Value
110.	Distance	I atti I I avei seu	(km)
1	TMC	TMC Airport – TMC Airport	(KIII)
1	Airport	Twe Anport – Twe Anport	0
2	Kita Beach	TMC Airport – Pantai Kita	19
2	Itita Deach	The Amport Tunun Nu	17
3	Kabisu	TMC Airport - Kabisu Lobo Oro	9
	Lobo Oro	Site	
	Site		
4	Lendongara	TMC Airport – Lendongara Hill	10
	Hill	· ·	
5	Waikelo	TMC Airport – Waikelo Beach	8.3
	Beach	*	
6	Kawona	TMC Airport – Kawona Beach	8.7
	Beach	-	
7	Karakat	TMC Airport – Karakat Indah	22
	Indah	Beach	
	Beach		
8	Waikuri	TMC Airport – Waikuri Lagoon	42
	Lagoon		
9	Mandorak	TMC Airport – Mandorak Beach	40
	Beach		
10	Tanjung	TMC Airport – Tanjung Karoso	48.7
	Karoso	Resort – Tanjung Karoso Beach	(from;
	Beach		51)
11	Tanjung	TMC Airport – Tanjung Karoso	43
	Karoso	Resort	
	Resort		
12	Rotenggaro	TMC Airport – Rotenggaro Beach	47
	Beach		
13	Bawana	TMC Airport – Bawana Beach	58
	Beach		
14	Watu	TMC Airport - Watu Malandong	63
	Malandong	Beach	
	Beach		
15	Sumba	TMC Airport – Sumba Culture	6.5
	Culture	House	
	House		
16	Praijing	TMC Airport – Praijing Village	43
	Village		

TABLE III. THE RESULT OF DIJKSTRA'S ALGORITHM

17	Lailiang Beach	TMC Airport – Praijing Village – Lailiang Beach	62 (from; 66)
18	Rua Beach	TMC Airport – Rua Beach	59
19	Mata Yangu Waterfall	TMC Airport – Mata Yangu Waterfall	65

Through the implementation of the program, it was found that Dijkstra's algorithm was able to generate a path from the starting point of TMC Airport to a particular location with a more optimal distance value, such as on the TMC Airport route to Tanjung Karoso Beach, the algorithm can provide route information that can be passed along the route. The total distance value is 48.7 km from the initial data of 51 km. Then, the following example on the TMC Airport path to Lailiang Beach obtained a route with a total distance value of 62 km, is smaller than the initial data of 66 km. Dijkstra's algorithm also computes other destinations until results are shown in Table III.

B. Bellman-Ford Algorithm

1) Calculation of the shortest distance

Calculations with the Bellman-Ford algorithm can be used to find the shortest distance between source nodes to each node. The Bellman-Ford algorithm can be used to find the shortest route solution of the all-pairs type, but in terms of implementation, the program is executed in the form of a single-source shortest path problem, namely from a certain node to all other nodes [24]. Calculation results are displayed with good and accurate results.

2) Complexity

For the complexity of the Bellman-Ford algorithm, it is considered quite complex with 409 lines of source code.

3) Execution Time

Based on 5 experiments on the Python program, the following data obtained for testing results on program execution as follows.

TABLE IV.	BELLMAN-FORD ALGORITHM EXECUTION TIME

Execution	Time
First Execution	0.0000004 seconds
Second Execution	0.0000009 second
Third Execution	0.0000005 seconds
Fourth Execution	0.0000006 seconds
Fifth Execution	0.0000004 seconds
Average Time	0.00000056 seconds

Results on testing show that programs with Bellman-Ford algorithms can run at very fast execution times, and can display fast execution results. Through these tests, the average execution time was 0.000056 seconds only.

4) Final Result

For the final results to be represented through rows and

columns, the following is a Table V for variable initialization for each tourist attraction based on the dataset used.

** • • •			a I			
Variable	Sumba	Variable	Sumba			
Initialization	Regional	Initialization	Regional			
	Destination		Destination			
1	TMC Airport	11	Tanjung Karoso			
	_		Resort			
2	Kita Beach	12	Rotenggaro			
			Beach			
3	Kabisu Lobo	13	Bawana Beach			
	Oro Site					
4	Lendongara Hill	14	Watu			
	-		Malandong			
			Beach			
5	Waikelo Beach	15	Sumba Culture			
			House			
6	Kawona Beach	16	Praijing Village			
7	Karakat Indah	17	Lailiang Beach			
	Beach					
8	Waikuri Lagoon	18	Rua Beach			
0	Mandanala	10	Mata Xanan			
9	Mandorak	19	Mata Yangu			
	Beach		Waterfall			
10	Tanjung Karoso					
	Beach					

Based on the table above, the final results in the Fig. 3 below are adjusted for the variable initialization in rows and columns with their respective distance values.

Vano	1	1.2	1	+	38	4	T			10.	11	10.1	11	14.	- 15	bah .	-17	18	111
1.	.0	19		.10.	10	8.7	22	42	- 10	48.7	-43	4T	18	61	-6.5	-41	102	- 59-	18
2.1	19	- 6	- 20	.11	- 22	36	-40	. 29	1.87	:447	10	43	78 .	:79	13	- 45	42	- 11	67
1		28	¥.	.11.	11	16	78	49	47	947	45	11	64	-89	13	49	18	- 85	11
4	.00	13	11	0	- 12	17	33.	. 60	- 19	85.7	38		68	20	14	- 90	15	:85	13
	83	12	12	15	8.	3.7	19	- 99	37	45	34.3	.90	017	85	11	47.	24	-83	78
	8.7	28	10	17	18.9	. 0.	14	3.4	34	11	39.3	47		+1	12	40	68	.64	71
+	22	4.8	30	11	100	14		21	1.0	22-	253	11	41	41	1.9	51	34	1.14	79
	-42	10		34	10401	14	26	0	3	34	4.1	- 21	11	100	101401	2.913	14	145	- 24
	.40	-12	47	40.	187	82	18	12		9.	3.6	18	28	34	-15-	317		10	. 87,5
18.	42.7	94.7	34.7	18.7	1461	141	27	-88		0	8.5	34	16	30	42.5	42.7	101	141	light -
11	-41-	49	-49		46.1	384	251	8.1	111	5.1		.14	18	35	1.7	77	- 94	46-	100
12	-41	: 65	53	34	50	45	31	: 20	18	39.	18	0	18	24	41	- 11	55	28	104
12	.76	19	- 54	66	117.5	58	-0	11	28	36	34	29		7.0	. 19	71	79.	. 69	1.94
14	45	1.79	40.	18	- 10	44	47.	28	194	30	.84	34	2.8	. 0	149	67	74	- 94	
11	6.5	D	12	14	11	UT.	11	1.87	100	42.7	.17	- 11	19	.97			10	17	64
18	-81	-49		3.8	1.87	49	37	315	1.94	#2.7	77	44	8.	NT.	1.46		19	23	24
11	42	1.42	- 64	69	- 64	14410	26	. 66	14	101	746	44	19	34	1.68	19	0.	20	1.46
18	.90	44	- 65	8.6	- 63	44	28.	. 50	12	11.	36	. 15	58	54	- 45	- 21	38		417
18	85	10	24	13	30	71	25	- 88	107	105	180	104	- 29	30		24	43	-40	-0

Fig. 3. The result of Bellman-Ford algorithm

The rows and columns in the table are sorted based on the initialization of the data that has been compiled. The Bellman-Ford algorithm works by carrying out calculations on each row and column and obtaining changes at several points with a smaller total distance traveled, changes are highlighted by the colored part. It was found that from all the existing data, the Bellman-Ford algorithm was able to analyze the data well and provide the shortest route solution correctly at each point.

So, by implementing the Bellman-Ford algorithm with a complexity of 409 lines of source code and several tests, it was found that the Bellman-Ford algorithm was considered effective in being able to find the minimum distance value of each node in the dataset. However, the efficiency of the program line is considered sufficient, because when compared to the other algorithms, the application of the source code to the Bellman-Ford algorithm is executed in large numbers.

C. Floyd-Warshall Algorithm

1) Calculation of the shortest distance

Calculations on the Floyd-Warshall algorithm are indeed more suitable for use for the all-pairs type [16], calculations are carried out using a Python program with 19 nodes which are then written in matrix form. From the implementation, can obtain that the calculations can run well and can find the most optimal solution, in this case the shortest route between all pairs of vertices.

2) Complexity

For the complexity of the program lines used, there are as many as 59 lines of source code.

3) Execution Time

Based on 5 experiments on the Python program, the following data obtained for testing results on Table VI program execution as follows.

TABLE VI.	FLOYD-WARSHALL ALGORITHM EXECUTION TIME

Execution	Time
First Execution	0.0000009 seconds
Second Execution	0.0000005 seconds
Third Execution	0.0000004 seconds
Fourth Execution	0.0000006 seconds
Fifth Execution	0.0000005 seconds
Average Time	0.00000058 seconds

Table VI depicts the results of each program execution process that applies the Floyd-Warshall algorithm. It can be seen that the execution time used to run the program is also very fast with the average execution time from 5 tests being 0.00000058 seconds.

4) Final Result

The final result of applying the algorithm based on each row and column is initialized as follows.

Variable Sumba Variable Sumba								
variable	Sumba	variable	Sumba					
Initialization	Regional	Initialization	Regional					
	Destination		Destination					
1	TMC Airport	11	Tanjung Karoso					
			Resort					
2	Kita Beach	12	Rotenggaro					
			Beach					
3	Kabisu Lobo	13	Bawana Beach					
	Oro Site							
4	Lendongara Hill	14	Watu					
			Malandong					
			Beach					
5	Waikelo Beach	15	Sumba Culture					
			House					
6	Kawona Beach	16	Praijing Village					
7	Karakat Indah	17	Lailiang Beach					
	Beach							
8	Waikuri Lagoon	18	Rua Beach					

9	Mandorak	19	Mata Yangu
	Beach		Waterfall
10	Tanjung Karoso		
	Beach		

The table above represents the variable initialization for each tourist attraction, then the results will be made in the form of rows and columns represented by the initialization as in Table VII. The following is a Fig. 4 for the results after applying the Floyd-Warshall algorithm to the dataset used.

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Fig. 4. The result of Floyd-Warshall algorithm

Rows and columns are also written based on the initialization of data that represents each location. Calculations are carried out at each point and obtained using the Floyd-Warshall algorithm calculation method which is also able to display route results with more optimal distance values at several points compared to the initial data, depicted in colored sections. It was found that this algorithm can provide quite significant changes to the problem of finding the shortest route.

So, by applying the Floyd-Warshall algorithm with a source code line complexity of 59 lines and several tests, it can be obtained that the Floyd-Warshall algorithm is considered effective and efficient in finding the minimum distance value for each node in the dataset. This algorithm is also very effective for use in the problem of determining the shortest route of the all-pairs type.

D. Johnson's Algorithm

1) Calculation of the shortest distance

The stages of the Johnson algorithm calculation steps use an approach in determining the shortest path of all pairs of paths where the destination vertices in a graph are among others [32], as follows.

(a) In the first step, add a vertex S to all path points on graph G by giving it a set of 0.

(b) Run the Bellman-Ford algorithm on G' with node S as the source by finding the minimum weight, then calculating the heuristic or h[v-1] for every number of nodes in graph G.

(c) When calculating h[...], then recalculate the edges of the graph using the formula: w(u, v) = w(u, v) + h[u] - h[v].

(d) If the iteration is complete, then delete the added S node and all weights are now positive on the graph. If the calculation iteration has not been completed, repeat step (c).(e) Run Dijkstra's shortest path algorithm for each node as the source and calculate the shortest route (u,v).

(f) If the iteration is complete, the shortest route u and v will

be found. If the calculation iteration has not been completed, repeat step (e).

2) Complexity

The complexity of source code when running Johnson's Algorithm program in Python is 86 lines.

3) Execution Time

The time to run the Johnson's algorithm in Python program implemented 5 times, so that it can better determine the amount of time the program is running and can take the average time. The following Table VIII is the result of the execution of the program.

TABLE VIII. JOHNSON'S ALGORITHM EXECUTION TIME						
Execution	Time					
First Execution	0.01395870000000768 seconds					
Second Execution	0.0163969000000016 seconds					
Third Execution	0.01337990000000333 seconds					
Fourth Execution	0.01944430000003523 seconds					
Fifth Execution	0.0164161000000071 seconds					
Average Time	0.0159191800000010988 seconds					

Table VIII explains the test results of running the program with the Johnson's algorithm and is obtained from the first execution trial to the last trial stage. The time required to run the program has a very slight time difference, the time used is only around 0.01 seconds at each stage. With this time, it can be seen that the execution time for this algorithm is also very fast and from the five experiments, the average time required to run the program and display the results is 0.0159191800000010988 seconds on the console.

4) Final Result

The results are managed by Johnson's algorithm which then finds the results of the shortest distance weights with those that have been modified from vertex-1 to vertex-19 to several other vertices. The results of the distance when running the Johnson's algorithm program with the Python programming language, found 72 shorter path data from each node that was originally searched for all destination nodes.

So, in the data Johnson's algorithm looks for the distance from vertex-1 and vertex-19 to all destination data, by obtaining 86 lines of program complexity and 0.0159191800000010988 program time in seconds. The results obtained found 72 shorter path data from each initial node that was searched for all destination nodes.

E. Ant Colony Algorithm

1) Calculation of the shortest distance

The steps for calculating the Ant Colony algorithm use an approach in determining the shortest path starting from a TMC Airport node to every one pair of destination nodes on a graph, including the following:

(a) The first step, set the value of Ant Colony optimization on the initial node. With conditions, nodes that have been visited

are included in the tabu list, so they will not be visited again. (b) After performing probabilistic node calculations, then select the next node to be assigned the Ant Colony value. (c) Move to the next node and the visited node becomes tabu

(c) Move to the next hode and the visited hode becomes tabu list.

(d) View all visited nodes, if not repeat step (b).

(e) Then it will record the length of the side or the distance of the route taken, and delete the entire tabu list.

(f) Determine the shortest route from the current nodes and update the pheromone, if the Ant Colony travels one full route back to the beginning.

(g) If the limit for the number of Ant Colony is the maximum iteration that has been reached and will be completed, thus determining the shortest route. If the maximum iteration has not been completed, repeat step (a).

2) Complexity

The complexity of source code when running the Python language of Ant Colony algorithm program is 84 lines.

3) Execution Time

When running the Python language program, the Ant Colony algorithm implemented 5 times, so that it can better determine the amount of time the program is running and can take the average time. In testing the execution of the Python language program, it can be seen in the Table IX below the results of testing the program.

Execution	Time
First Execution	1.242826700000233 seconds
Second Execution	0.802237999999885 seconds
Third Execution	0.8272481999999854 seconds
Fourth Execution	1.2401753999999983 seconds
Fifth Execution	0.853923300000195 seconds
Average Time	0.993282320000003 seconds

TABLE IX. ANT COLONY ALGORITHM EXECUTION TIME

In contrast to several algorithms that have been tested previously, the execution time for running the Ant Colony program actually varies. From Table IX it can be seen that in the first and fourth experiments, the time used to run the program was more than 1 second. However, it is different for other experiments, which only take less than 1 second. Even so, each stage still has a slight time difference and is still considered very fast to be able to execute an Ant Colony program according to the program specifications used. Also obtained was the average program execution time of 0.993282320000003 seconds.

4) Final Result

In this algorithm, tests were carried out on 19 ants from 19 data on the destinations in the Sumba area used. Then on the 19 data used variable initialization is applied, as in the following Table X.

TABLE X. VARIABLE INITIALIZATION							
Variable Initialization	Sumba Regional	Variable Initialization	Sumba Regional				
muanzation	Destination	IIItiaiizatioii	Destination				
1	TMC Airport	11	Tanjung Karoso Resort				
2	Kita Beach	12	Rotenggaro Beach				
3	Kabisu Lobo Oro Site	13	Bawana Beach				
4	Lendongara Hill	14	Watu Malandong Beach				
5	Waikelo Beach	15	Sumba Culture House				
6	Kawona Beach	16	Praijing Village				
7	Karakat Indah Beach	17	Lailiang Beach				
8	Waikuri Lagoon	18	Rua Beach				
9	Mandorak Beach	19	Mata Yangu Waterfall				
10	Tanjung Karoso Beach						

Then, testing was carried out and the shortest distance was found from the iteration of the ant in the form of a circuit. The following Table XI represents the test results data with the Ant Colony algorithm.

TABLE XI. ANT COLONY ALGORITHM RESULTS

Ant	Path Traversed	Value
Iteration		(km)
1 st	1 - 15 - 3 - 4 - 2 - 16 - 17 - 18 - 19 - 13 - 14	385.3
	- 12 - 10 - 8 - 9 - 11 - 7 - 6 - 5 - 1	
2 nd	1 - 15 - 3 - 2 - 4 - 5 - 6 - 7 - 8 - 9 - 11 -	342.0
	10 - 12 - 13 - 14 - 18 - 17 - 16 - 19 - 1	
3 rd	1 - 15 - 5 - 6 - 7 - 9 - 8 - 11 - 10 - 12 - 14	376.0
	- 13 - 17 - 16 - 19 - 18 - 2 - 4 - 3 - 1	
4^{th}	1 - 15 - 5 - 6 - 7 - 9 - 8 - 11 - 10 - 12 - 13	382.0
	- 14 - 19 - 16 - 18 - 17 - 3 - 4 - 2 - 1	
5 th	1 - 15 - 6 - 5 - 3 - 4 - 2 - 18 - 17 - 16 - 19	402.0
-th	- 13 - 14 - 10 - 11 - 12 - 8 - 9 - 7 - 1	120.0
6 ^m	1 - 6 - 5 - 3 - 4 - 2 - 16 - 17 - 18 - 19 - 7 -	428.0
7th	15 - 9 - 8 - 11 - 10 - 12 - 13 - 14 - 1	274.0
/	1 - 6 - 5 - 15 - 3 - 4 - 2 - 16 - 19 - 18 - 1/	374.0
oth	-15 - 14 - 12 - 10 - 11 - 9 - 8 - 7 - 1	280.0
0	- 9 - 8 - 10 - 11 - 15 - 4 - 2 - 3 - 1	380.0
9 th	1 - 15 - 6 - 5 - 3 - 4 - 2 - 11 - 10 - 9 - 8 -	395.0
-	7 - 12 - 13 - 14 - 18 - 17 - 16 - 19 - 1	07010
10 th	1 - 3 - 4 - 2 - 5 - 6 - 7 - 8 - 9 - 11 - 10 -	337.5
	12 - 13 - 14 - 18 - 17 - 16 - 19 - 15 - 1	
11 th	1 - 15 - 3 - 4 - 2 - 6 - 5 - 9 - 10 - 11 - 12 -	452.0
	13 - 14 - 17 - 18 - 16 - 19 - 7 - 8 - 1	
12 th	1 - 5 - 6 - 15 - 3 - 4 - 2 - 16 - 19 - 7 - 8 -	447.0
	9 - 11 - 10 - 12 - 14 - 13 - 17 - 18 - 1	
13 th	1 - 15 - 3 - 4 - 2 - 5 - 6 - 7 - 9 - 8 - 10 -	380.0
1.4th	11 - 12 - 13 - 14 - 16 - 17 - 18 - 19 - 1	242.0
14"	1 - 15 - 3 - 5 - 6 - 7 - 9 - 8 - 11 - 10 - 12 -	343.0
15th	14 - 13 - 16 - 17 - 10 - 19 - 2 - 4 - 1 $1 - 15 - 6 - 5 - 7 - 8 - 9 - 11 - 10 - 12 - 14$	380.0
15	- 13 - 19 - 16 - 17 - 18 - 3 - 4 - 2 - 1	389.0
16 th	1 - 5 - 6 - 7 - 8 - 9 - 11 - 10 - 12 - 14 - 13	385.0
	- 19 - 18 - 17 - 16 - 2 - 15 - 3 - 4 - 1	20210
17 th	1 - 15 - 7 - 6 - 5 - 3 - 4 - 2 - 18 - 17 - 16 -	400.0
	19 - 13 - 14 - 12 - 10 - 11 - 9 - 8 - 1	
18 th	1 - 15 - 5 - 6 - 7 - 9 - 8 - 11 - 10 - 12 - 13	353.0

	- 14 - 18 - 17 - 19 - 16 - 3 - 4 - 2 - 1	
19 th	1 - 5 - 6 - 15 - 7 - 8 - 9 - 11 - 10 - 12 - 13 - 14 - 18 - 16 - 17 - 19 - 3 - 4 - 2 - 1	377.0

The display of running iterations is accompanied by path traversed, and how many total distance values are taken on the corresponding route. In this Ant Colony data algorithm, find the distance from TMC Airport by visiting all 18 destinations and then returning to TMC Airport, obtaining 84 lines of program complexity and 0.993282320000003 program time in seconds. The most optimal route results obtained are as far as 337.5 km, with the following routes: TMC Airport - Kabisu Lobo Oro Site - Lendongara Hill - Kita Beach - Waikelo Beach - Kawona Beach - Karakat Indah Beach - Waikuri Lagoon - Mandorak Beach - Tanjung Karoso Resort - Tanjung Beach Karoso - Rotenggaro Beach - Bawana Beach - Watu Malandong Beach - Rua Beach -Lailiang Beach - Praijing Village - Mata Yangu Waterfall -Sumba Cultural House - TMC Airport.

F. Discussion

Through the results described, the analysis performed on the performance of each algorithm on the test. Through testing algorithms judged based on the parameters that were indicators of this study, each algorithm had its own advantages and disadvantages. As such, the Dijkstra's algorithm can provide the most optimal shortest route results, but it can only effectively work to solve single-source and single-destination-type problems. The Bellman-Ford algorithm also has its own advantages, such as being able to hold negative weight values on graphs, but in this study it has the most complexity on source code among other algorithms. The Bellman-Ford algorithm can be used for almost the same path as Dijkstra in single-source and single-destination types. The Floyd-Warshall algorithm is more effective for use in allpairs path types, and can provide the most optimal value at any point. Johnson's algorithm is a combination of Dijkstra's algorithm and Bellman-Ford's algorithm. Meanwhile, for the Ant Colony algorithm, it works with the longest execution time among other algorithms, but after analysis it turns out that the Ant Colony algorithm has a high degree of effectiveness in resolving the problem of route lookup for all nodes.

Speaking of which algorithms are best suited to implement, it is also necessary to see how the graph picture will be examined and which solutions will be sought for the problem. Thus, the search can be performed more effectively and efficiently and have the most optimal distance value, in the sense that it does not focus only on which route with the least weight value. As described, Dijkstra's and Bellman-Ford algorithms are suitable for use in single-source and single-destination graphs, Floyd-Warshall algorithms for allpairs type, Johnson's algorithm can also be used for all-pairs type searches and Ant Colony can be used to find routes that can visit all destinations at once at a time well and the minimum possible distance value.

IV. CONCLUSION

Based on testing of the five algorithms and adjusted to the parameters that have been compiled, it is found that each

algorithm has its own type. Dijkstra's algorithm can be used to calculate the shortest route for single-source and singledestination types. It's the same as the Bellman-Ford algorithm, except that the Bellman-Ford algorithm can be used at the same time on graphs that have negative weight values. Meanwhile, the Floyd-Warshall algorithm is suitable for use on the all-pairs type. For Johnson's algorithm it can be used to determine the shortest path from all pairs of paths where the destination node is on a graph, and Ant Colony for calculating from a node to every one pair of destination nodes. Through the implementation of the Python program, it was found that there was a change in the dataset used, namely the Southwest Sumba data, to become data with a more optimum distance value, in this case the minimum distance value from a starting point to a destination point.

Through the Python program that has been researched, the complexity of each of the various algorithms is obtained. Dijkstra's algorithm has 70 lines of source code, Bellman-Ford has 409 source codes, Floyd-Warshall has 59 source codes, Johnson has 86 source codes, and Ant Colony has 84 source codes. Thus, the algorithm with the lowest complexity is owned by Floyd-Warshall algorithm.

Then, Dijkstra's algorithm with execution time running in the Python program obtained an average time of 0.00064035999999873252 seconds, Bellman-Ford's algorithm for 0.00000056 seconds, Floyd-Warshall's algorithm for 0.00000058 seconds, Johnson's algorithm for 0.0159191800000010988 seconds, and Ant Colony's algorithm for 0.993282320000003 seconds. Thus, the Bellman-Ford and Floyd-Warshall algorithms are considered to have the fastest execution time.

Based on all the parameters that have been examined, it can be concluded that the Ant Colony algorithm has high effectiveness for solving the most efficient route finding for all nodes. And the Floyd-Warshall algorithm is considered to be the most effective algorithm in finding all-pairs type routes, and is able to find the most optimal distance value between one node and another node. Also, Bellman-Ford has the fastest execution time when a Python program is run.

Through such implementations it is also obtained about how the algorithm performs if it is based through certain parameters. Each algorithm has its own advantages and disadvantages when implemented. The journal focuses on how to test the subject, then provides a new perspective on which algorithms have the most efficient performance with good accuracy when implemented on specific graph models and path models.

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